

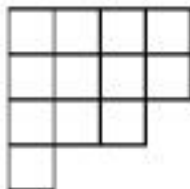
OK

**partition**  $\lambda \vdash n$ :  $\lambda = (\lambda_1, \lambda_2, \dots)$

$$\lambda_1 \geq \lambda_2 \geq \dots \geq 0$$

$$\sum \lambda_i = n$$

**(Young) diagram** of  $\lambda = (4, 4, 3, 1)$ :

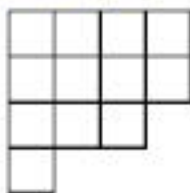


**partition**  $\lambda \vdash n$ :  $\lambda = (\lambda_1, \lambda_2, \dots)$

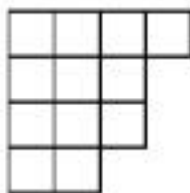
$$\lambda_1 \geq \lambda_2 \geq \dots \geq 0$$

$$\sum \lambda_i = n$$

**(Young) diagram** of  $\lambda = (4, 4, 3, 1)$ :



Young diagram of the **conjugate** partition  $\lambda' = (4, 3, 3, 2)$ :



standard Young tableau (SYT) of  
shape  $\lambda \vdash n$ , e.g.,  $\lambda = (4, 4, 3, 1)$ :

<

	1	2	7	10
	3	5	8	12
^	4	6	11	
	9			

standard Young tableau (SYT) of  
 shape  $\lambda \vdash n$ , e.g.,  $\lambda = (4, 4, 3, 1)$ :

<

	1	2	7	10
	3	5	8	12
^	4	6	11	
	9			

$f^\lambda = \#$  of SYT of shape  $\lambda$

E.g.,  $f^{(3,2)} = 5$ :

123	124	125	134	135
45	35	34	25	24

$\exists$  simple formula for  $f^\lambda$  (Frame-Robinson-  
 Thrall **hook-length formula**)

**Note.**  $f^\lambda = \dim(\text{irrep. of } \mathfrak{S}_n)$ , where  $\mathfrak{S}_n$  is the symmetric group of all permutations of  $1, 2, \dots, n$ .

**Note.**  $f^\lambda = \dim(\text{irrep. of } \mathfrak{S}_n)$ , where  $\mathfrak{S}_n$  is the **symmetric group** of all permutations of  $1, 2, \dots, n$ .

**RSK algorithm:** a bijection

$$w \xrightarrow{\text{rsk}} (P, Q),$$

where  $w \in \mathfrak{S}_n$  and  $P, Q$  are SYT of the same shape  $\lambda \vdash n$ .

Write  $\lambda = \mathbf{sh}(w)$ , the **shape** of  $w$ .

**Note.**  $f^\lambda = \dim(\text{irrep. of } \mathfrak{S}_n)$ , where  $\mathfrak{S}_n$  is the **symmetric group** of all permutations of  $1, 2, \dots, n$ .

**RSK algorithm:** a bijection

$$w \xrightarrow{\text{RSK}} (P, Q),$$

where  $w \in \mathfrak{S}_n$  and  $P, Q$  are SYT of the same shape  $\lambda \vdash n$ .

Write  $\lambda = \mathbf{sh}(w)$ , the **shape** of  $w$ .

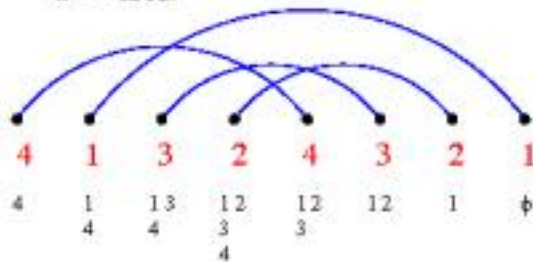
**R** = Gilbert de Beauregard Robinson

**S** = Craige Schensted (= Ed. Eds)

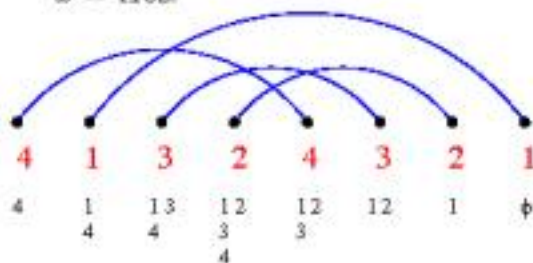
**K** = Donald Ervin Knuth



$w = 4132$ :



$w = 4132$ :



$$(P, Q) = \begin{pmatrix} 12 & 13 \\ 3 & 2 \\ 4 & 4 \end{pmatrix}$$

**Schensted's theorem:** Let  $w \xrightarrow{\text{rank}}$   
 $(P, Q)$ , where  $sh(P) = sh(Q) = \lambda$ .  
Then

$$is(w) = \text{longest row length} = \lambda_1$$

$$ds(w) = \text{longest column length} = \lambda'_1.$$

**Summing and averaging dimensions**

$$\begin{aligned} \dim \text{span}(v_1, \dots, v_n) &= \dim \langle v_1, \dots, v_n \rangle \\ &= \dim \text{row}(A) \\ &= \dim \text{row}(A) + \dim \text{null}(A) \\ &= \dim(A) \end{aligned}$$

**Application: optimal encoding**

**Example:** messages have to be sent to  $n$  receivers  $v_1, \dots, v_n$  in the network  $G$ . Each message takes one time unit to be sent along an edge of  $G$ .



**Def:** Send every message  $i$ .

**Definition:** send every message  $i$  at least  $d_i$  times.

**Flow conditions:**

- 1. The outgoing flow to each node must be the same.
- 2. There is no flow out of the source, the sink, the  $v_i$ .

problem: find flow  $f = (f_e)$

$$\begin{aligned} f_e + f_{e'} + \dots &= d_i \\ \sum_e f_e &= 0 \end{aligned}$$

(Formal) degree of the  $(v_i, v_j) \in E$



Using degree of the receiver  $v_i$  problem  $f = (f_e)$



instead: Using volume (WF) of degree  $d_i$  for  $v_i$  for  $(v_i, v_j) \in E$



$$f^* = \# \text{ of WF of degree } d_i$$

Ex:  $d_i = 4$

01	11	01	11	01
00	10	00	10	00

2 nodes found for  $f^*$  (Flow Condition) both have degree 4

**Def:**  $f^*$  is the flow of  $(G, d_i)$  when  $v_i$  is the receiver,  $v_j$  is the source,  $v_k$  is the sink,  $v_l$  is the source.

**Def:** algorithm to find  $f^*$

$$= \# \langle f \rangle$$

where  $\langle f \rangle$  is the set of WF of the same degree  $d_i$ .

What is  $\# \langle f \rangle$  the class of  $f$ ?

- 1. is  $\#$  of WF of degree  $d_i$
- 2. is  $\#$  of WF of degree  $d_i$
- 3. is  $\#$  of WF of degree  $d_i$

**Schensted's theorem:** Let  $w \xrightarrow{\text{rak}}$   
 $(P, Q)$ , where  $sh(P) = sh(Q) = \lambda$ .  
Then

$$is(w) = \text{longest row length} = \lambda_1$$

$$ds(w) = \text{longest column length} = \lambda'_1.$$

$$4132 \xrightarrow{\text{rak}} \begin{pmatrix} 12 & 13 \\ 3 & 2 \\ 4 & 4 \end{pmatrix}$$

$$is(w) = 2, \quad ds(w) = 3$$

**Schensted's theorem:** Let  $w \xrightarrow{\text{rank}} (P, Q)$ , where  $sh(P) = sh(Q) = \lambda$ .  
Then

$$is(w) = \text{longest row length} = \lambda_1$$

$$ds(w) = \text{longest column length} = \lambda'_1.$$

**Corollary** (Erdős-Szekeres, Seidenberg). Let  $w \in \mathfrak{S}_{pq+1}$ . Then either  $is(w) > p$  or  $ds(w) > q$ .

**Proof.** Let  $\lambda = sh(w)$ . If  $is(w) \leq p$  and  $ds(w) \leq q$  then  $\lambda_i \leq p$  and  $\lambda'_i \leq q$ , so  $\sum \lambda_i \leq pq$ .  $\square$

**Corollary.** Say  $p \leq q$ . Then

$$\begin{aligned} \#\{w \in \mathfrak{S}_{pq} : \text{is}(w) = p, \text{ds}(w) = q\} \\ = \left(f(p^q)\right)^2 \end{aligned}$$

By hook-length formula, this is

$$\left( \frac{(pq)!}{1!2^2 \cdots p^p(p+1)^{p-1} \cdots q^q(q+1)^{q-1} \cdots (p+q-1)!} \right)^2$$

**Romik:** let

$$w \in \mathfrak{S}_{p^2}, \quad \text{is}(w) = \text{ds}(w) = p.$$

Let  $P_w$  be the permutation matrix of  $w$  with corners  $(\pm 1, \pm 1)$ . Then (informally) as  $p \rightarrow \infty$  almost surely the 1's in  $P_w$  will become dense in the region bounded by the curve

$$(x^2 - y^2)^2 + 2(x^2 + y^2) = 3,$$

and will remain isolated outside this region.



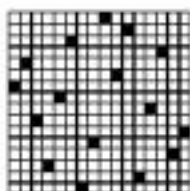
**Romik:** let

$$w \in \mathfrak{S}_{p^2}, \quad \text{is}(w) = \text{ds}(w) = p.$$

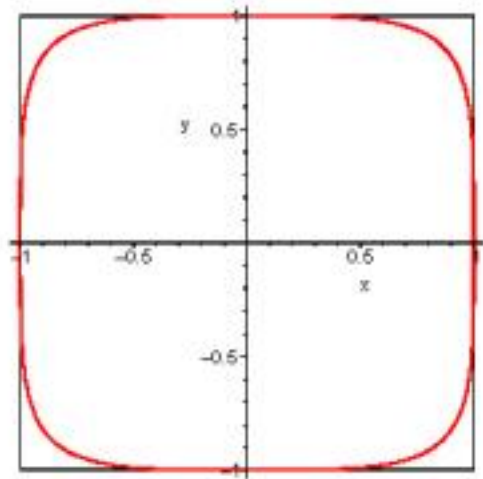
Let  $P_w$  be the permutation matrix of  $w$  with corners  $(\pm 1, \pm 1)$ . Then (informally) as  $p \rightarrow \infty$  almost surely the 1's in  $P_w$  will become dense in the region bounded by the curve

$$(x^2 - y^2)^2 + 2(x^2 + y^2) = 3,$$

and will remain isolated outside this region.



$$w = 9, 11, 6, 14, 2, 10, 1, 5, 13, 3, 16, 8, 15, 4, 12, 7$$



$$(x^2 - y^2)^2 + 2(x^2 + y^2) = 3$$

## Distribution of $is(w)$

$$\begin{aligned} E(n) &= \text{expectation of } is(w), w \in \mathfrak{S}_n \\ &= \frac{1}{n!} \sum_{\lambda \vdash n} \lambda_1 (f^\lambda)^2 \end{aligned}$$

## Distribution of $is(w)$

$$\begin{aligned} E(n) &= \text{expectation of } is(w), w \in \mathfrak{S}_n \\ &= \frac{1}{n!} \sum_{\lambda \vdash n} \lambda_1 (f^\lambda)^2 \end{aligned}$$

**Ulam:** what is distribution of  $is(w)$ ?  
rate of growth of  $E(n)$ ?

## Distribution of $is(w)$

$$\begin{aligned} E(n) &= \text{expectation of } is(w), w \in \mathfrak{S}_n \\ &= \frac{1}{n!} \sum_{\lambda \vdash n} \lambda_1 (f^\lambda)^2 \end{aligned}$$

**Ulam:** what is distribution of  $is(w)$ ?  
rate of growth of  $E(n)$ ?

**Hammersley** (1972):

$$\exists c = \lim_{n \rightarrow \infty} \frac{E(n)}{\sqrt{n}},$$

and

$$\frac{\pi}{2} \leq c \leq e.$$

Conjectured  $c = 2$ .

Logan-Shepp, Vershik-Kerov (1977):

$c = 2$

Idea of proof.

$$E(n) = \frac{1}{n!} \sum_{\lambda \vdash n} \lambda_1 (f^\lambda)^2$$
$$\approx \frac{1}{n!} \max_{\lambda \vdash n} \lambda_1 (f^\lambda)^2.$$

Find "limiting shape" of  $\lambda \vdash n$  maximizing  $\lambda$  as  $n \rightarrow \infty$  using hook-length formula.

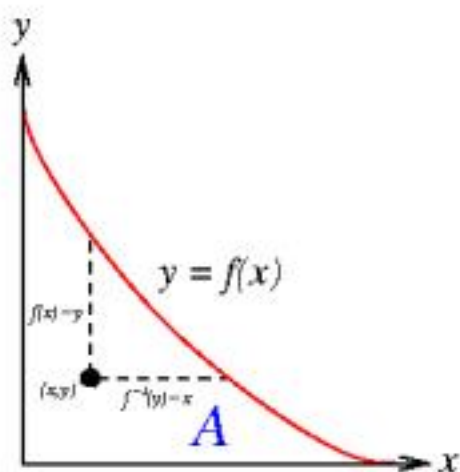
## Increasing and decreasing sub-sequences

3 18**4**9**6**725 (i.s.)

31**8**4967**2**5 (d.s.)

$is(w) = |\text{longest i.s.}| = 4$

$ds(w) = |\text{longest d.s.}| = 3$

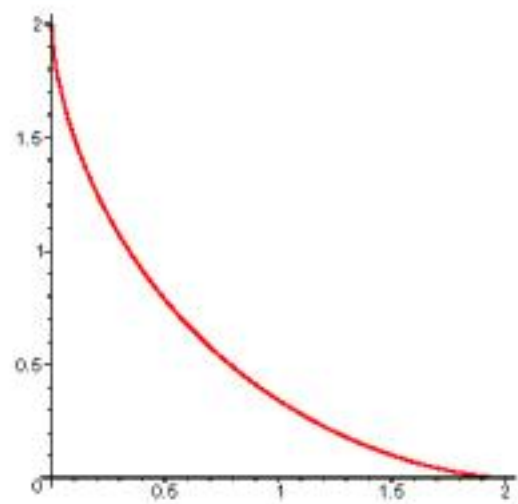


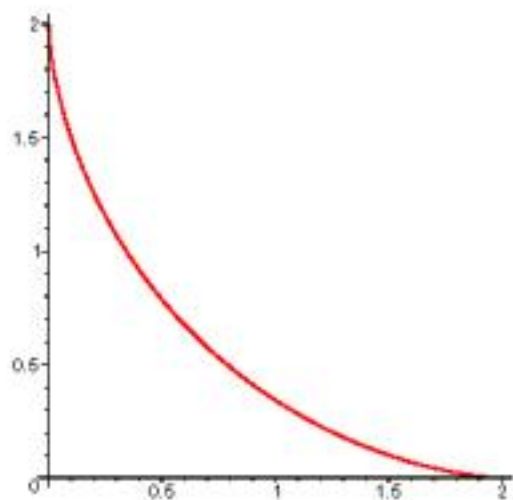
$$\min \iint_A \log(f(x) + f^{-1}(y) - x - y) dx dy,$$

subject to

$$\iint_A dx dy = 1.$$







$$x = y + 2 \cos \theta$$

$$y = \frac{2}{\pi} (\sin \theta - \theta \cos \theta)$$

$$0 \leq \theta \leq \pi$$

$$u_k(n) := \#\{w \in \mathfrak{S}_n : \text{is}_n(w) \leq k\}.$$

$u_k(n) := \#\{w \in \mathfrak{S}_n : \text{is}_n(w) \leq k\}$ .

J. M. Hammersley (1972):

$$u_2(n) = C_n = \frac{1}{n+1} \binom{2n}{n},$$

is Catalan number.

$u_k(n) := \#\{w \in \mathfrak{S}_n : \text{is}_n(w) \leq k\}$ .

J. M. Hammersley (1972):

$$u_2(n) = C_n = \frac{1}{n+1} \binom{2n}{n},$$

a Catalan number.

For  $\geq 130$  combinatorial interpretations  
of  $C_n$ , see

[www-math.mit.edu/~rstan/ec](http://www-math.mit.edu/~rstan/ec)

I. Gessel (1990):

$$\sum_{n \geq 0} w_k(n) \frac{x^{2n}}{n! 2^n} = \det \left[ I_{|i-j|}(2x) \right]_{i,j=1}^k,$$

where

$$I_m(2x) = \sum_{j \geq 0} \frac{x^{m+2j}}{j!(m+j)!},$$

a **hyperbolic Bessel function** of the first kind of order  $m$ .

I. Gessel (1990):

$$\sum_{n \geq 0} u_k(n) \frac{x^{2n}}{n!^2} = \det \left[ I_{|i-j|}(2x) \right]_{i,j=1}^k,$$

where

$$I_m(2x) = \sum_{j \geq 0} \frac{x^{m+2j}}{j!(m+j)!},$$

a **hyperbolic Bessel function** of the first kind of order  $m$ .

E. g.,

$$\begin{aligned} \sum_{n \geq 0} u_2(n) \frac{x^{2n}}{n!^2} &= U_0(2x)^2 - U_1(2x)^2 \\ &= \sum_{n \geq 0} C_n \frac{x^{2n}}{n!^2}. \end{aligned}$$

**Corollary.** For fixed  $k$ ,  $u_k(n)$  is **P-recursive**, e.g.,

$$\begin{aligned} & (n+4)(n+3)^2 u_4(n) \\ = & (20n^3 + 62n^2 + 22n - 24)u_4(n-1) \\ & - 64n(n-1)^2 u_4(n-2) \end{aligned}$$

$$\begin{aligned} & (n+6)^2(n+4)^2 u_6(n) \\ = & (375 - 400n - 843n^2 - 322n^3 - 35n^4)u_6(n-1) \\ & + (259n^2 + 622n + 45)(n-1)^2 u_6(n-2) \\ & - 225(n-1)^2(n-2)^2 u_6(n-3). \end{aligned}$$

Conjectures on form of recurrence due to Bergeron, Favreau, and Krob.



### Baik-Deift-Johansson:

Define  $u(x)$  by

$$\frac{d^2}{dx^2}u(x) = 2u(x)^3 + xu(x) \quad (*),$$

with certain initial conditions.

(\*) is the **Painlevé II** equation (roughly, the branch points and essential singularities are independent of the initial conditions).

**Application:** airplane boarding

**Naive model:** passengers board in order  $w = a_1 a_2 \cdots a_n$  for seats  $1, 2, \dots, n$ . Each passenger takes one time unit to be seated after arriving at his seat.



## Paul Painlevé

**1863:** born in Paris.

## Paul Painlevé

**1863:** born in Paris.

**1890:** Grand Prix des Sciences Mathématiques

## Paul Painlevé

**1863:** born in Paris.

**1890:** Grand Prix des Sciences Mathématiques

**1908:** first passenger of Wilbur Wright;  
set flight duration record of one hour, 10  
minutes.

## Paul Painlevé

**1863:** born in Paris.

**1890:** Grand Prix des Sciences Mathématiques

**1908:** first passenger of Wilbur Wright;  
set flight duration record of one hour, 10  
minutes.

**1917, 1925:** Prime Minister of France.

## Paul Painlevé

**1863:** born in Paris.

**1890:** Grand Prix des Sciences Mathématiques

**1908:** first passenger of Wilbur Wright;  
set flight duration record of one hour, 10  
minutes.

**1917, 1925:** Prime Minister of France.

**1933:** died in Paris.

Tracy-Widom distribution:

$$F(t) = \exp\left(-\int_t^\infty (x-t)u(x)^2 dx\right)$$



Tracy-Widom distribution:

$$F(t) = \exp\left(-\int_t^\infty (x-t)u(x)^2 dx\right)$$

**Theorem** (Baik-Deift-Johansson) For random (uniform)  $w \in \mathfrak{S}_n$  and all  $t \in \mathbb{R}$  we have

$$\lim_{n \rightarrow \infty} \text{Prob}\left(\frac{\text{is}_n(w) - 2\sqrt{n}}{n^{1/6}} \leq t\right) = F(t).$$

Tracy-Widom distribution:

$$F(t) = \exp\left(-\int_t^\infty (x-t)u(x)^2 dx\right)$$

**Theorem** (Baik-Deift-Johansson) For random (uniform)  $w \in \mathfrak{S}_n$  and all  $t \in \mathbb{R}$  we have

$$\lim_{n \rightarrow \infty} \text{Prob}\left(\frac{\text{is}_n(w) - 2\sqrt{n}}{n^{1/6}} \leq t\right) = F(t).$$

**Corollary.**

$$\begin{aligned} \text{is}_n(w) &= 2\sqrt{n} + \left(\int t dF(t)\right) n^{1/6} + o(n^{1/6}) \\ &= 2\sqrt{n} - (1.7711 \dots) n^{1/6} + o(n^{1/6}) \end{aligned}$$

Gessel's theorem reduces the problem to "just" analysis, viz, the **Riemann-Hilbert problem** in the theory of integrable systems, and the **method of steepest descent** to analyze the asymptotic behavior of integrable systems.

Gessel's theorem reduces the problem to "just" analysis, viz, the **Riemann-Hilbert problem** in the theory of integrable systems, and the **method of steepest descent** to analyze the asymptotic behavior of integrable systems.

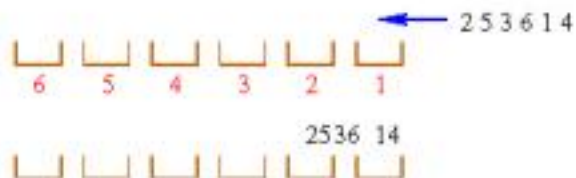
Where did the Tracy-Widom distribution  $F(t)$  come from?

---

$$F(t) = \exp\left(-\int_t^\infty (x-t)u(x)^2 dx\right)$$
$$\frac{d^2}{dx^2}u(x) = 2u(x)^3 + xu(x) \quad (*),$$

**Application:** airplane boarding

**Naive model:** passengers board in order  $w = a_1 a_2 \cdots a_n$  for seats  $1, 2, \dots, n$ . Each passenger takes one time unit to be seated after arriving at his seat.



### Gaussian Unitary Ensemble (GUE):

Consider an  $n \times n$  hermitian matrix  $M = (M_{ij})$  with probability density

$$Z_n^{-1} e^{-\text{tr}(M^2)} dM,$$

$$dM = \prod_i dM_{ii} \cdot \prod_{i < j} d(\text{Re}(M_{ij})) d(\text{Im}(M_{ij})),$$

where  $Z_n$  is a normalization constant.

**Tracy-Widom** (1994): let  $\alpha_1$  denote the largest eigenvalue of  $M$ . Then

$$\lim_{n \rightarrow \infty} \text{Prob} \left( \left( \alpha_1 - \sqrt{2n} \right) \sqrt{2n}^{1/6} \leq t \right) = F(t).$$

Is the connection between  $is(w)$  and  
GUE a coincidence?



Is the connection between  $is(w)$  and GUE a coincidence?

Okounkov provides a connection, via the theory of **random topologies on surfaces**. Very briefly, a surface can be described in two ways:

- Gluing polygons along their edges, connected to random matrices via quantum gravity.
- Ramified covering of a sphere, which can be formulated in terms of permutations.

Joint with:

Bill Chen 陈永川

Eva Deng 邓玉平

Rosena Du 杜若霞

Catherine Yan 颜华菲

(complete) matching:



**crossing:** 

**nesting:** 

(complete) matching:



crossing:



nesting:



total number of matchings on  $[2n] :=$   
 $\{1, 2, \dots, 2n\}$  is

$$(2n - 1)!! := 1 \cdot 3 \cdot 5 \cdots (2n - 1).$$

(complete) matching:



crossing:



nesting:



total number of matchings on  $[2n] := \{1, 2, \dots, 2n\}$  is

$$(2n - 1)!! := 1 \cdot 3 \cdot 5 \cdots (2n - 1).$$

**Theorem.** *The number of matchings on  $[2n]$  with no crossings (or with no nestings) is*

$$C_n := \frac{1}{n+1} \binom{2n}{n}.$$

Well-known:

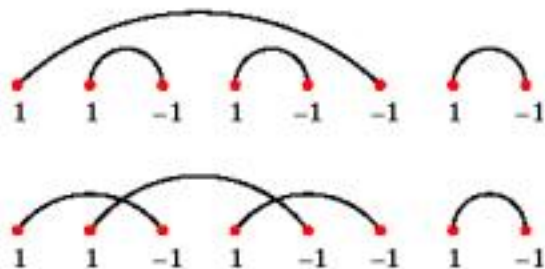
$$C_n = \#\{a_1 \cdots a_{2n} : a_i = \pm 1, \\ a_1 + \cdots + a_i \geq 0, \sum a_i = 0\}$$

(ballot sequence).

Well-known:

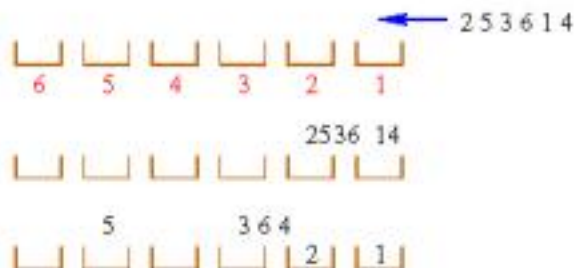
$$C_n = \#\{a_1 \cdots a_{2n} : a_i = \pm 1, \\ a_1 + \cdots + a_i \geq 0, \sum a_i = 0\}$$

(ballot sequence).



**Application:** airplane boarding

**Naive model:** passengers board in order  $w = a_1 a_2 \cdots a_n$  for seats  $1, 2, \dots, n$ . Each passenger takes one time unit to be seated after arriving at his seat.





What is the analogue of increasing and decreasing subsequences for matchings  $M$ ?

What is the analogue of increasing and decreasing subsequences for matchings  $M$ ?

Associate with a matching  $M$  on the vertices  $1, 2, \dots, 2n$  a fixed-point free involution  $w_M \in \mathfrak{S}_{2n}$ :



$$w_M = (1, 5)(2, 7)(3, 4)(6, 8)$$

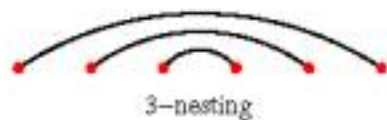
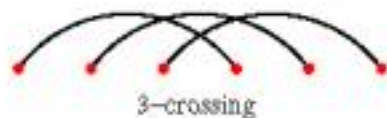
What is the analogue of increasing and decreasing subsequences for matchings  $M$ ?

Associate with a matching  $M$  on the vertices  $1, 2, \dots, 2n$  a fixed-point free involution  $w_M \in \mathfrak{S}_{2n}$ :



$$w_M = (1, 5)(2, 7)(3, 4)(6, 8)$$

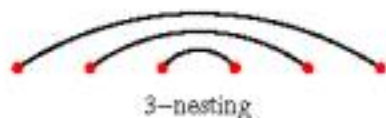
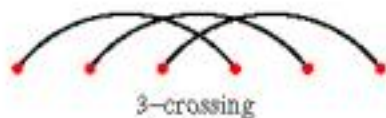
**Flaw:** no symmetry between is and ds (different distributions on fixed-point free involutions).



$M$  = matching

$$\text{cr}(M) = \max\{k : \exists k\text{-crossing}\}$$

$$\text{ne}(M) = \max\{k : \exists k\text{-nesting}\} = \frac{1}{2} \text{ds}(w_M)$$



$M$  = matching

$\text{cr}(M) = \max\{k : \exists k\text{-crossing}\}$

$\text{ne}(M) = \max\{k : \exists k\text{-nesting}\} = \frac{1}{2} \text{ds}(w_M)$

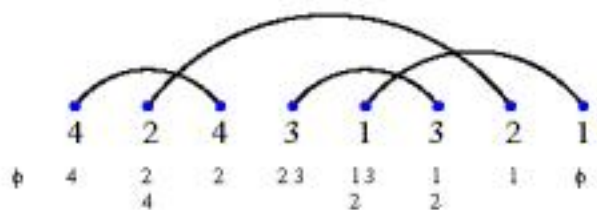
**Theorem.** Let  $f_n(i, j) = \#$  matchings  $M$  on  $[2n]$  with  $\text{cr}(M) = i$  and  $\text{ne}(M) = j$ . Then  $f_n(i, j) = f_n(j, i)$ .

**Corollary.**  $\#$  matchings  $M$  on  $[2n]$  with  $\text{cr}(M) = k$  equals  $\#$  matchings  $M$  on  $[2n]$  with  $\text{ne}(M) = k$ .

Main tool: oscillating tableaux.

$\phi$

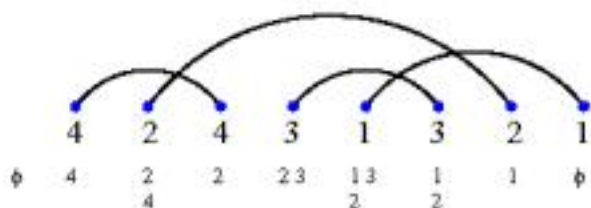
shape (3, 1), length 8



Main tool: oscillating tableaux.

$\phi$   $\square$   $\begin{smallmatrix} \square \\ \square \end{smallmatrix}$   $\square$   $\square\square$   $\begin{smallmatrix} \square & \square \\ \square & \square \end{smallmatrix}$   $\square\square$   $\square\square\square$   $\begin{smallmatrix} \square & \square \\ \square & \square \end{smallmatrix}$

shape (3, 1), length 8



$\Phi(M) = (\phi \square \begin{smallmatrix} \square \\ \square \end{smallmatrix} \square \square\square \begin{smallmatrix} \square & \square \\ \square & \square \end{smallmatrix} \square\square \square \phi)$

$\Phi$  is a bijection from matchings on  $1, 2, \dots, 2n$  to oscillating tableaux of length  $2n$ , shape  $\emptyset$ .



$\Phi$  is a bijection from matchings on  $1, 2, \dots, 2n$  to oscillating tableaux of length  $2n$ , shape  $\emptyset$ .

**Corollary.** *Number of oscillating tableaux of length  $2n$ , shape  $\emptyset$ , is  $(2n-1)!!$  (related to **Brauer algebra** of dimension  $(2n-1)!!$ ).*

$\Phi$  is a bijection from matchings on  $1, 2, \dots, 2n$  to oscillating tableaux of length  $2n$ , shape  $\emptyset$ .

**Corollary.** *Number of oscillating tableaux of length  $2n$ , shape  $\emptyset$ , is  $(2n-1)!!$  (related to **Brauer algebra** of dimension  $(2n-1)!!$ ).*

**Schensted's theorem for matchings.** *Let*

$$\Phi(M) = (\emptyset = \lambda^0, \lambda^1, \dots, \lambda^{2n} = \emptyset).$$

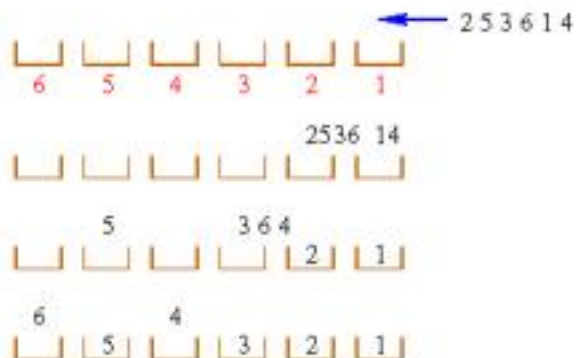
*Then*

$$\begin{aligned} \text{cr}(M) &= \max\{(\lambda^i)_1' : 0 \leq i \leq n\} \\ \text{ne}(M) &= \max\{\lambda_1^i : 0 \leq i \leq n\}. \end{aligned}$$

**Proof.** Reduce to ordinary RSK.

**Application:** airplane boarding

**Naive model:** passengers board in order  $w = a_1 a_2 \cdots a_n$  for seats  $1, 2, \dots, n$ . Each passenger takes one time unit to be seated after arriving at his seat.



**Easy:** Total waiting time =  $is(w)$ .

Now let  $cr(M) = i$ ,  $ne(M) = j$ , and

$$\Phi(M) = (\emptyset = \lambda^0, \lambda^1, \dots, \lambda^{2n} = \emptyset).$$

Define  $M'$  by

$$\Phi(M') = (\emptyset = (\lambda^0)', (\lambda^1)', \dots, (\lambda^{2n})' = \emptyset).$$

By Schensted's theorem for matchings,

$$cr(M') = j, \quad ne(M') = i.$$

Thus  $M \mapsto M'$  is an involution on matchings of  $[2n]$  interchanging  $cr$  and  $ne$ .

Now let  $cr(M) = i$ ,  $ne(M) = j$ , and

$$\Phi(M) = (\emptyset = \lambda^0, \lambda^1, \dots, \lambda^{2n} = \emptyset).$$

Define  $M'$  by

$$\Phi(M') = (\emptyset = (\lambda^0)', (\lambda^1)', \dots, (\lambda^{2n})' = \emptyset).$$

By Schensted's theorem for matchings,

$$cr(M') = j, \quad ne(M') = i.$$

Thus  $M \mapsto M'$  is an involution on matchings of  $[2n]$  interchanging  $cr$  and  $ne$ .

$\Rightarrow$  **Theorem.** Let  $f_n(i, j) = \#$  matchings  $M$  on  $[2n]$  with  $cr(M) = i$  and  $ne(M) = j$ . Then  $f_n(i, j) = f_n(j, i)$ .

Now let  $\text{cr}(M) = i$ ,  $\text{ne}(M) = j$ , and

$$\Phi(M) = (\emptyset = \lambda^0, \lambda^1, \dots, \lambda^{2n} = \emptyset).$$

Define  $M'$  by

$$\Phi(M') = (\emptyset = (\lambda^0)', (\lambda^1)', \dots, (\lambda^{2n})' = \emptyset).$$

By Schensted's theorem for matchings,

$$\text{cr}(M') = j, \quad \text{ne}(M') = i.$$

Thus  $M \mapsto M'$  is an involution on matchings of  $[2n]$  interchanging  $\text{cr}$  and  $\text{ne}$ .

$\Rightarrow$  **Theorem.** Let  $f_n(i, j) = \#$  matchings  $M$  on  $[2n]$  with  $\text{cr}(M) = i$  and  $\text{ne}(M) = j$ . Then  $f_n(i, j) = f_n(j, i)$ .

**Open:** simple description of  $M \mapsto M'$ , the analogue of

$$a_1 a_2 \cdots a_n \mapsto a_n \cdots a_2 a_1,$$

which interchanges  $is$  and  $ds$

### Enumeration of $k$ -noncrossing matchings (or nestings).

**Recall:** The number of matchings  $M$  on  $[2n]$  with no crossings, i.e.,  $\text{cr}(M) = 1$ , (or with no nestings) is  $C_n = \frac{1}{n+1} \binom{2n}{n}$ .

What about the number with  $\text{cr}(M) \leq k$ ?

Assume  $\text{cr}(M) \leq k$ . Let

$$\Phi(M) = (\emptyset = \lambda^0, \lambda^1, \dots, \lambda^{2n} = \emptyset).$$

Regard each  $\lambda^i = (\lambda_1^i, \dots, \lambda_k^i) \in \mathbb{N}^k$ .

**Corollary.** The number  $f_k(n)$  of matchings  $M$  on  $[2n]$  with  $\text{cr}(M) \leq k$  is the number of lattice paths of length  $2n$  from  $\mathbf{0}$  to  $\mathbf{0}$  in the region

$$\mathcal{C}_n := \{(a_1, \dots, a_k) \in \mathbb{N}^k : a_1 \leq \dots \leq a_k\}$$

with steps  $\pm e_i$  ( $e_i = i$ th unit coordinate vector).

$\mathcal{C}_n \otimes \mathbb{R}_{\geq 0}$  is a fundamental chamber for the Weyl group of type  $B_k$ .



**Grabiner-Magyar:** applied **Gessel-Zeilberger reflection principle** to solve this lattice path problem (not knowing connection with matchings).

**Grabiner-Magyar:** applied **Gessel-Zeilberger reflection principle** to solve this lattice path problem (not knowing connection with matchings).

**Theorem.** Define

$$H_k(x) = \sum_n f_k(n) \frac{x^{2n}}{(2n)!}.$$

Then

$$H_k(x) = \det \left[ I_{|i-j|}(2x) - I_{i+j}(2x) \right]_{i,j=1}^k$$

where

$$I_m(2x) = \sum_{j \geq 0} \frac{x^{m+2j}}{j!(m+j)!}$$

as before.

**Example.**  $k = 1$  (noncrossing matchings):

$$\begin{aligned} H_1(x) &= I_0(2x) - I_2(2x) \\ &= \sum_{j \geq 0} C_j \frac{x^{2j}}{(2j)!}. \end{aligned}$$

**Example.**  $k = 1$  (noncrossing matchings):

$$\begin{aligned}H_1(x) &= I_0(2x) - I_2(2x) \\ &= \sum_{j \geq 0} C_j \frac{x^{2j}}{(2j)!}.\end{aligned}$$

**Compare:**

$u_k(n) := \#\{w \in \mathfrak{S}_n : \text{longest increasing subsequence of length } \leq k\}.$

$$\sum_{n \geq 0} u_k(n) \frac{x^{2n}}{n!} = \det [I_{i-j}(2x)]_{i,j=1}^k.$$

**Baik-Rains** (implicitly):

$$\lim_{n \rightarrow \infty} \text{Prob} \left( \frac{\text{cr}_n(M) - \sqrt{2n}}{(2n)^{1/6}} \leq \frac{t}{2} \right) = F_1(t),$$

where

$$F_1(t) = \sqrt{F(t)} \exp \left( \frac{1}{2} \int_t^\infty u(x) dx \right),$$

where  $F(t)$  is the Tracy-Widom distribution and  $u(x)$  the Painlevé II function.

---

$$= \exp \left( - \int_t^\infty (x-t) u(x)^2 dx \right)$$

$$\frac{d^2}{dx^2} u(x) = 2u(x)^3 + xu(x)$$

Bachmat, et al.: more sophisticated model

Two conclusions:

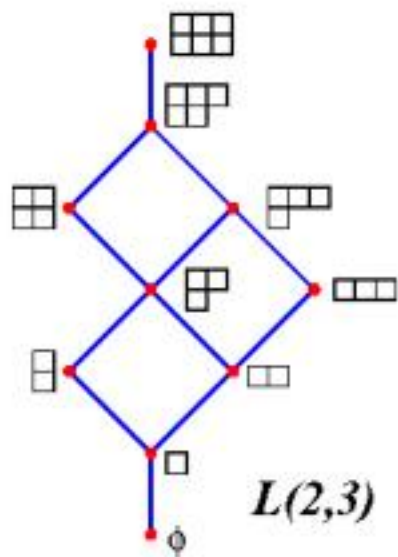
- Usual system (back-to-front) not much better than random.
- Better: first board window seats, then center, then aisle.

$$g_{j,k}(n) := \#\{\text{matchings } M \text{ on } [2n], \\ \text{cr}(M) \leq j, \text{ne}(M) \leq k\}$$

Now

$g_{j,k}(n) = \#\{(\emptyset = \lambda^0, \lambda^1, \dots, \lambda^{2n} = \emptyset) :$   
 $\lambda^{i+1} = \lambda^i \pm \square, \lambda^i \subseteq j \times k \text{ rectangle}\},$   
 a walk on the Hasse diagram  $\mathcal{H}(j, k)$   
 of

$L(j, k) := \{\lambda \subseteq j \times k \text{ rectangle}\},$   
 ordered by inclusion.





$\mathbf{A}$  = adjacency matrix of  $\mathcal{H}(j, k)$

$\mathbf{A}_0$  = adjacency matrix of  $\mathcal{H}(j, k) - \{\emptyset\}$ .

Transfer-matrix method  $\Rightarrow$

$$\sum_{n \geq 0} g_{j,k}(n) x^{2n} = \frac{\det(I - xA_0)}{\det(I - xA)}.$$

$\mathbf{A}$  = adjacency matrix of  $\mathcal{H}(j, k)$

$\mathbf{A}_0$  = adjacency matrix of  $\mathcal{H}(j, k) - \{\emptyset\}$ .

Transfer-matrix method  $\Rightarrow$

$$\sum_{n \geq 0} g_{j,k}(n) x^{2n} = \frac{\det(I - xA_0)}{\det(I - xA)}.$$

**Theorem** (Grabiner, implicitly) Every zero of  $\det(I - xA)$  has the form

$$2(\cos(\pi r_1/m) + \cdots + \cos(\pi r_j/m)),$$

where each  $r_i \in \mathbb{Z}$  and  $m = j + k + 1$ .

$\mathbf{A}$  = adjacency matrix of  $\mathcal{H}(j, k)$

$\mathbf{A}_0$  = adjacency matrix of  $\mathcal{H}(j, k) - \{\emptyset\}$ .

Transfer-matrix method  $\Rightarrow$

$$\sum_{n \geq 0} g_{j,k}(n) x^{2n} = \frac{\det(I - xA_0)}{\det(I - xA)}.$$

**Theorem** (Grabiner, implicitly) Every zero of  $\det(I - xA)$  has the form

$$2(\cos(\pi r_1/m) + \cdots + \cos(\pi r_j/m)),$$

where each  $r_i \in \mathbb{Z}$  and  $m = j + k + 1$ .

**Corollary.** Every factor of  $\det(I - xA)$  over  $\mathbb{Q}$  has degree dividing

$$\frac{1}{2}\phi(2(j+k+1)),$$

where  $\phi$  is the Euler phi-function.

**Example.**

$$j = 2, k = 5, \frac{1}{2}\phi(16) = 4:$$

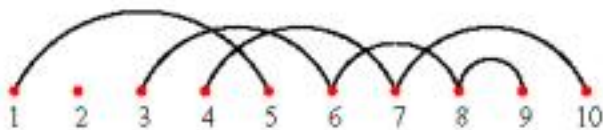
$$\begin{aligned} \det(I - xA) &= (1 - 2x^2)(1 - 4x^2 + 2x^4) \\ &= (1 - 8x^2 + 8x^4)(1 - 8x^2 + 8x^3 - 2x^4) \\ &= (1 - 8x^2 - 8x^3 - 2x^4) \end{aligned}$$

$$j = k = 3, \frac{1}{2}\phi(14) = 3:$$

$$\begin{aligned} \det(I - xA) &= (1 - x)(1 + x)(1 + x - 9x^2 - x^3) \\ &= (1 - x - 9x^2 + x^3)(1 - x - 2x^2 + x^3)^2 \\ &= (1 + x - 2x^2 - x^3)^2 \end{aligned}$$

**Partition** of the set  $[n]$ :

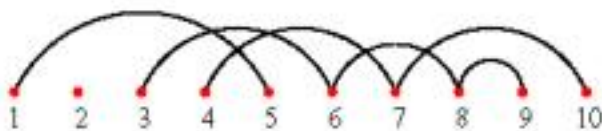
$\{1, 5\}, \{2\}, \{3, 6, 8, 9\}, \{4, 7, 10\}$



Generalize oscillating tableaux to  
**vacillating tableaux** (related to the  
**partition algebra**).

**Partition** of the set  $[n]$ :

$\{1, 5\}, \{2\}, \{3, 6, 8, 9\}, \{4, 7, 10\}$



Generalize oscillating tableaux to  
**vacillating tableaux** (related to the  
**partition algebra**).

Many other variations: see paper!

**partition**  $\lambda \vdash n$ :  $\lambda = (\lambda_1, \lambda_2, \dots)$

$$\lambda_1 \geq \lambda_2 \geq \dots \geq 0$$

$$\sum \lambda_i = n$$